

BAR-AND-JOINT RIGIDITY ON THE MOMENT CURVE COINCIDES WITH COFACTOR RIGIDITY ON A CONIC

Luis Crespo Ruiz^{*1} and Francisco Santos^{*2}

^{1,2}*Department of Mathematics, Statistics and Computing, Universidad de Cantabria, Santander, Spain*
luis.cresporuiz@unican.es , francisco.santos@unican.es

Submitted: Oct 5, 2021; Accepted: Jan 4, 2023; Published: Mar 15, 2023

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Abstract. We show that, for points along the moment curve, the bar-and-joint rigidity matroid and the hyperconnectivity matroid coincide, and that both coincide with the C_{d-1}^{d-2} -cofactor rigidity of points along any (non-degenerate) conic in the plane. For hyperconnectivity in dimension two, having the points in the moment curve is no loss of generality.

We also show that, restricted to bipartite graphs, the bar-and-joint rigidity matroid is freer than the hyperconnectivity matroid.

Keywords. Rigidity, hyperconnectivity, moment curve, cofactor rigidity

Mathematics Subject Classifications. 52C25, 52B40

1. Introduction

Rigidity matroids were introduced by Graver in [Gra91] as a way to formalize and generalize rigidity theory. Besides the matroids associated to (infinitesimal) *bar-and-joint rigidity* they include two other important examples: the matroids of C_{d-1}^{d-2} -*cofactor rigidity* introduced by Whiteley [Whi96] and of the *hyperconnectivity* introduced by Kalai [Kal85]. See the precise definitions below.

^{*}Supported by grants PID2019-106188GB-I00 funded by MCIN/AEI/10.13039/501100011033 and by FPU19/04163 of the Spanish Government and by project CLaPPo (21.SI03.64658) of Universidad de Cantabria and Banco Santander.

We show an interesting case in which the three theories are equivalent:

Theorem 1.1. *Let d be a positive integer and let $t_1, \dots, t_n \in \mathbb{R}$ be (distinct) real numbers. Then, the following three d -dimensional rigidity matroids coincide:*

- *The C_{d-1}^{d-2} -cofactor matroid of points $\{(t_i, t_i^2)\}_{i=1}^n$ along the standard parabola.*
- *The bar-and-joint matroid of points $\{(t_i, \dots, t_i^d)\}_{i=1}^n$ along the moment curve.*
- *The d -hyperconnectivity matroid of vectors $\{(1, t_i, \dots, t_i^{d-1})\}_{i=1}^n$, which, if no t_i equals zero, is equivalent to that of $\{(t_i, \dots, t_i^d)\}_{i=1}^n$.*

The *moment curve* $\{(t, \dots, t^d) : t \in \mathbb{R}\} \subset \mathbb{R}^d$ is an important curve from several points of view. From the perspective of algebraic geometry, it is the rational normal curve of degree d (in a specific embedding, but any other embedding is algebraically isomorphic to it). From the perspective of polytope theory, any finite subset of n points in the moment curve defines a *cyclic polytope*, an example of a *neighborly polytope*: a d -polytope with n vertices that attains the maximum possible number of faces of each dimension. From the perspective of rigidity theory, one important property of the moment curve is that it is contained in many (linearly independent) quadrics. By a result of Bolker and Roth [BR80] this has drastic consequences for the bar-and-joint rigidity of bipartite graphs embedded with vertices in it. See Lemma 4.7 and Remark 4.8.

For $d = 2$ the moment curve is the standard parabola. Since the three forms of rigidity are projectively invariant, any non-degenerate conic can be substituted for the standard parabola, for arbitrary d in part (1), or for $d = 2$ in parts (2) and (3).

Let us recall some of the needed rigidity concepts. Good comprehensive references for background are [NSW21, Whi96]; [CJT22] also contains everything we need. For hyperconnectivity see [Kal85] or [JT21].

A *point configuration* is a finite list of distinct points $\mathbf{p} = (\mathbf{p}_1, \dots, \mathbf{p}_n)$ in \mathbb{R}^d . A *framework* on \mathbf{p} is a simple graph with vertex set $\{\mathbf{p}_1, \dots, \mathbf{p}_n\}$. Since the points in \mathbf{p} are labelled by $[n]$, we identify frameworks on \mathbf{p} with subsets of $\binom{[n]}{2}$. The *bar-and-joint (infinitesimal) rigidity matrix* of \mathbf{p} is the following $\binom{n}{2} \times nd$ matrix:

$$R(\mathbf{p}) := \begin{pmatrix} \mathbf{p}_1 - \mathbf{p}_2 & \mathbf{p}_2 - \mathbf{p}_1 & 0 & \dots & 0 & 0 \\ \mathbf{p}_1 - \mathbf{p}_3 & 0 & \mathbf{p}_3 - \mathbf{p}_1 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & & \vdots & \vdots \\ \mathbf{p}_1 - \mathbf{p}_n & 0 & 0 & \dots & 0 & \mathbf{p}_n - \mathbf{p}_1 \\ 0 & \mathbf{p}_2 - \mathbf{p}_3 & \mathbf{p}_3 - \mathbf{p}_2 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & & \vdots & \vdots \\ 0 & 0 & 0 & \dots & \mathbf{p}_{n-1} - \mathbf{p}_n & \mathbf{p}_n - \mathbf{p}_{n-1} \end{pmatrix}$$

That is to say, each row r_{ij} of $R(\mathbf{p})$ ($\{i, j\} \in \binom{[n]}{2}$) consists of n blocks of size d and has zeroes everywhere except in the i th and j th blocks, where it has, respectively, $\mathbf{p}_i - \mathbf{p}_j$ and $\mathbf{p}_j - \mathbf{p}_i$.

Definition 1.2. The *bar-and-joint rigidity matroid* of \mathbf{p} is the linear matroid of rows of $R(\mathbf{p})$, with ground set $\binom{[n]}{2}$. We denote it $\mathcal{R}_d(\mathbf{p})$. A framework on \mathbf{p} is called *rigid* (respectively *stress-free*, *isostatic*) if the corresponding set of rows is spanning, (respectively *independent*, *a basis*) in the rigidity matroid.

The reason for this terminology is that, if $R(\mathbf{p})|_G$ denotes the row-submatrix of $R(\mathbf{p})$ corresponding to a framework G , then the kernel of $R(\mathbf{p})|_G$ are the infinitesimal motions of the points (“joints”) that preserve lengths of all edges in G (“bars”); the kernel of its transpose are the equilibrium stresses on G : assignments of extension/contraction forces to the edges that cancel out at every vertex, so that the system is in equilibrium.

We can also give up symmetry and directly define the following *hyperconnectivity matrix* of \mathbf{p} , introduced by Kalai [Kal85]:

$$H(\mathbf{p}) := \begin{pmatrix} \mathbf{p}_2 & -\mathbf{p}_1 & 0 & \dots & 0 & 0 \\ \mathbf{p}_3 & 0 & -\mathbf{p}_1 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & & \vdots & \vdots \\ \mathbf{p}_n & 0 & 0 & \dots & 0 & -\mathbf{p}_1 \\ 0 & \mathbf{p}_3 & -\mathbf{p}_2 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & & \vdots & \vdots \\ 0 & 0 & 0 & \dots & \mathbf{p}_n & -\mathbf{p}_{n-1} \end{pmatrix}.$$

Definition 1.3. We call *hyperconnectivity matroid* of \mathbf{p} , and denote it $\mathcal{H}_d(\mathbf{p})$, the linear matroid of rows of $H(\mathbf{p})$.

For even d , $\mathcal{H}_d(n)$ coincides with the algebraic matroid of skew-symmetric matrices of rank at most d , of interest in low-rank matrix completion problems. See, e.g., [Ber17]. This relation is well-known, but we offer a proof in Proposition 3.1. In the case $d = 2$ this is, in turn, the same as the algebraic matroid of the Plücker embedding of the Grassmannian $Gr(2, \mathbb{R})$ in $\mathbb{R}^{\binom{n}{2}}$. The following combinatorial characterization of $\mathcal{H}_2(n)$ is known [Ber17]: a graph $G \subset \binom{[n]}{2}$ is independent in $\mathcal{H}_2(n)$ if, and only if, G admits an acyclic orientation with no alternating closed walk.

Finally, for a point configuration $\mathbf{p} \subset \mathbb{R}^2$ in the plane and a positive integer d we define the *cofactor rigidity matrix* of degree $d - 1$,

$$C_d(\mathbf{p}) := \begin{pmatrix} c(\mathbf{p}_1 - \mathbf{p}_2) & -c(\mathbf{p}_1 - \mathbf{p}_2) & 0 & \dots & 0 & 0 \\ c(\mathbf{p}_1 - \mathbf{p}_3) & 0 & -c(\mathbf{p}_1 - \mathbf{p}_3) & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & & \vdots & \vdots \\ c(\mathbf{p}_1 - \mathbf{p}_n) & 0 & 0 & \dots & 0 & -c(\mathbf{p}_1 - \mathbf{p}_n) \\ 0 & c(\mathbf{p}_2 - \mathbf{p}_3) & -c(\mathbf{p}_2 - \mathbf{p}_3) & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & & \vdots & \vdots \\ 0 & 0 & 0 & \dots & c(\mathbf{p}_{n-1} - \mathbf{p}_n) & -c(\mathbf{p}_{n-1} - \mathbf{p}_n) \end{pmatrix},$$

where $c(x, y) = (x^{d-1}, x^{d-2}y, \dots, xy^{d-2}, y^{d-1})$.

Definition 1.4. We call *cofactor rigidity matroid* of \mathbf{p} , and denote it $\mathcal{C}_{d-1}^{d-2}(\mathbf{p})$, the linear matroid of rows of $C_d(\mathbf{p})$.

This formalism was introduced by Whiteley in the 1990's [Whi96], elaborating on previous work of Billera on smooth bivariate splines [Bil88]. The notation \mathcal{C}_{d-1}^{d-2} comes from the relation of this theory to C^{d-2} -splines of degree $d-1$.

The three rigidity theories share important properties. Let \mathbf{p} and \mathbf{q} be two point configurations of size $n \geq d$ in general position in \mathbb{R}^d and \mathbb{R}^2 respectively.¹ Then:

1. $\text{rank}(\mathcal{R}_d(\mathbf{p})) = \text{rank}(\mathcal{H}_d(\mathbf{p})) = \text{rank}(\mathcal{C}_{d-1}^{d-2}(\mathbf{q})) = nd - \binom{d+1}{2}$.
2. Every copy of the complete graph K_{d+2} is a circuit in $\mathcal{R}_d(\mathbf{p})$, $\mathcal{H}_d(\mathbf{p})$ and $\mathcal{C}_{d-1}^{d-2}(\mathbf{q})$. (And hence every copy of K_m , with $m \leq d+1$, is independent).

By a theorem of Nguyen [Ngu10], matroids on the ground set $\binom{[n]}{2}$ and satisfying these two properties coincide exactly with the *abstract rigidity matroids of dimension d* introduced by Graver [Gra91].

For reasons that will become apparent in the proof, we call the matroid in Theorem 1.1 the *polynomial rigidity matroid of degree $d-1$ with parameters t_1, \dots, t_n* (Definition 2.4). We denote it $\mathcal{P}_d(t_1, \dots, t_n)$. Since points on the moment curve are in general position, $\mathcal{P}_d(t_1, \dots, t_n)$ is an abstract rigidity matroid too.

There is some interest among the experts on the relations between the three matroids $\mathcal{H}_d(\mathbf{p})$, $\mathcal{R}_d(\mathbf{p})$, and $\mathcal{C}_{d-1}^{d-2}(\mathbf{q})$ in the case when \mathbf{p} and \mathbf{q} are generic. We denote these matroids $\mathcal{H}_d(n)$, $\mathcal{R}_d(n)$, and $\mathcal{C}_{d-1}^{d-2}(n)$, and denote by $\mathcal{P}_d(n)$ the generic case of $\mathcal{P}_d(t_1, \dots, t_n)$.

It is conjectured that $\mathcal{C}_{d-1}^{d-2}(n)$ is freer than $\mathcal{R}_d(n)$ and that the latter is freer than $\mathcal{H}_d(n)$ (Conjecture 4.3). It is also conjectured that $\mathcal{C}_{d-1}^{d-2}(n)$ is the freest abstract rigidity matroid and that $\mathcal{H}_d(n)$ is the freest matroid in which every K_{d+2} and every $K_{d+1, d+1}$ are circuits (Conjectures 4.1 and 4.2). We recall what is known about these conjectures in Section 4, including a proof of the following result which is, to the best of our knowledge, new: restricted to bipartite graphs, $\mathcal{H}_d(n)$ coincides with $\mathcal{R}_d(\mathbf{p})$ where \mathbf{p} has the points of the two parts lying in two hyperplanes, but otherwise generic. In particular, on bipartite graphs $\mathcal{R}_d(n)$ is freer than $\mathcal{H}_d(n)$ (Theorem 4.4 and Corollary 4.5).

We hope that our new matroid $\mathcal{P}_d(n)$ may help shed light on these conjectures, since it shows a case where the three rigidity theories coincide. We finish Section 4 and the paper proving several properties that are common to $\mathcal{H}_d(n)$ and $\mathcal{P}_d(n)$ (Corollary 4.9 and Propositions 4.10 and 4.11). This suggests that perhaps $\mathcal{H}_d(n) = \mathcal{P}_d(n)$ (Question 4.6). For $d=2$ this is true, by the invariance of $\mathcal{H}_d(n)$ under linear scaling of the positions of points (Theorem 2.2).

2. Proof of Theorem 1.1

The matroids \mathcal{C}_{d-1}^{d-2} and \mathcal{R}_d are invariant under projective transformation in $\mathbb{RP}^d \supset \mathbb{R}^d$, see e.g. [NSW21]. \mathcal{H}_d is, as far as we know, only invariant under linear transformation or, rather, under projective transformation in \mathbb{RP}^{d-1} as a quotient space of $\mathbb{R}^d \setminus \{0\}$:

¹For \mathcal{R}_d and \mathcal{C}_{d-1}^{d-2} we here mean *affine* general position; that is, all affinely dependent subsets are affinely spanning. For \mathcal{H}_d we mean *linear* general position: all linearly dependent subsets are linearly spanning.

Lemma 2.1. Let $\mathbf{p} = (\mathbf{p}_1, \dots, \mathbf{p}_n)$ and $\mathbf{q} = (\mathbf{q}_1, \dots, \mathbf{q}_n)$ be point configurations in \mathbb{R}^d and assume that \mathbf{q} is obtained from \mathbf{p} by one of the following two procedures:

1. A global linear transformation $l : \mathbb{R}^d \rightarrow \mathbb{R}^d$ with $l(\mathbf{p}_i) = \mathbf{q}_i$ for all i .
2. Multiplication of each \mathbf{p}_i by a non-zero scalar α_i , so that $\mathbf{q}_i = \alpha_i \mathbf{p}_i$ for each i .

Then, $\mathcal{H}_d(\mathbf{p}) = \mathcal{H}_d(\mathbf{q})$.

Proof. In the first case, the matrix $H(\mathbf{q})$ is obtained from $H(\mathbf{p})$ multiplying on the right by a block-diagonal matrix having (the matrix of) l in each block. In the second case, it is obtained by first multiplying each row (i, j) by $\alpha_i \alpha_j$ and then dividing each column in block i by α_i . Multiplying rows or columns by non-zero scalars does not change the matroid. \square

This implies that \mathcal{H}_2 coincides with \mathcal{P}_2 :

Theorem 2.2. Let $\mathbf{p} = (\mathbf{p}_1, \dots, \mathbf{p}_n)$ be points in $(\mathbb{R} \setminus \{0\})^2$, with $\mathbf{p}_i = (a_i, b_i)$. Let $t_i = b_i/a_i$ for $i = 1, \dots, n$. Then, $\mathcal{H}_2(\mathbf{p})$ equals $\mathcal{P}_2(t_1, \dots, t_n)$. In particular, $\mathcal{H}_2(n) = \mathcal{P}_2(n)$.

Proof. Apply Lemma 2.1 with $\alpha_i = b_i/a_i^2$. \square

Let $\mathbb{R}[t]_{<d}$ denote the d -dimensional vector space of univariate polynomials of degree less than d with real coefficients. For each choice of parameters t_1, \dots, t_n and choice of one basis $F^i(t) = (f_1^i(t), \dots, f_d^i(t))$ of $\mathbb{R}[t]_{<d}$ for each i , construct the following *polynomial rigidity matrix of degree $d - 1$* :

$$P_{F^1, \dots, F^n}(t_1, \dots, t_n) := \begin{pmatrix} F^1(t_2) & -F^2(t_1) & 0 & \dots & 0 & 0 \\ F^1(t_3) & 0 & -F^3(t_1) & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & & \vdots & \vdots \\ F^1(t_n) & 0 & 0 & \dots & 0 & -F^n(t_1) \\ 0 & F^2(t_3) & -F^3(t_2) & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & & \vdots & \vdots \\ 0 & 0 & 0 & \dots & F^{n-1}(t_n) & -F^n(t_{n-1}) \end{pmatrix}.$$

Lemma 2.3. The linear matroid of rows of $P_{F^1, \dots, F^n}(t_1, \dots, t_n)$ is independent of the choice of bases F^1, \dots, F^n .

Proof. Given two choices of bases (F^1, \dots, F^n) and (G^1, \dots, G^n) , the matrix corresponding to G can be obtained from that of F multiplying on the right by the $nd \times nd$ block-diagonal matrix that has in the i -th diagonal block the $d \times d$ matrix that changes from basis F^i to basis G^i . \square

Definition 2.4. We denote this matroid $\mathcal{P}_d(t_1, \dots, t_n)$ and call it the *polynomial rigidity matroid of degree $d - 1$ with parameters (t_1, \dots, t_n)* .

Theorem 1.1 follows from the following more precise statement:

Theorem 2.5. The three matroids of Theorem 1.1 coincide with $\mathcal{P}_d(t_1, \dots, t_n)$.

Proof. Consider the following bases of $\mathbb{R}[t]_{<d}$, depending on t_i :

$$F^{t_i}(t) = \left(\frac{t_i^k - t^k}{t_i - t} \right)_{k=1, \dots, d}, \quad G^{t_i}(t) = ((t_i + t)^{k-1})_{k=1, \dots, d}.$$

Then, $P_{F^{t_1}, \dots, F^{t_n}}(t_1, \dots, t_n)$ equals the bar-and-joint rigidity matrix of points along the moment curve, except each row (i, j) has been divided by $t_i - t_j$, which does not affect the matroid. Similarly, $P_{G^{t_1}, \dots, G^{t_n}}(t_1, \dots, t_n)$ equals the cofactor rigidity matrix of points along the parabola, except each row (i, j) has been divided by $(t_i - t_j)^{d-1}$. For the latter, observe that for points $(x_i, y_i) = (t_i, t_i^2)$ along the parabola we have

$$(x_i - x_j)^{d-k} (y_i - y_j)^{k-1} = (t_i - t_j)^{d-k} (t_i^2 - t_j^2)^{k-1} = (t_i - t_j)^{d-1} (t_i + t_j)^{k-1}.$$

Finally, the hyperconnectivity matrix of $\{(1, \dots, t_i^{d-1})\}_{i \in [n]}$ equals $P_{F, \dots, F}(t_1, \dots, t_n)$ where we choose all bases equal to the monomial basis $F = (1, t, t^2, \dots, t^{d-1})$. \square

3. Hyperconnectivity and low rank skew-symmetric matrices

Recall that the *algebraic matroid* of an irreducible variety $V \subset \mathbb{R}^N$ is the matroid with ground set N and in which a subset of coordinates is independent if there is no (non-trivial) polynomial relation among them on V . Put differently, X is independent if $I(V) \cap \mathbb{R}[X] = 0$. When V is parametrized by a polynomial map $T : \mathbb{R}^M \rightarrow V \subset \mathbb{R}^N$, the algebraic matroid of V equals the linear matroid of rows of the Jacobian of T at a sufficiently generic point of \mathbb{R}^M [Ros14, Proposition 2.5]. (It is important that we are looking at algebraic varieties over \mathbb{R} . Algebraic matroids over finite fields are not all representable as linear matroids over those same fields).

We are interested in the variety $\mathcal{S}_d(n) \subset \mathbb{R}^{\binom{n}{2}}$ of skew-symmetric matrices of rank (at most) d , where each matrix is represented as the list of its entries with $1 \leq i < j \leq n$. We denote $\mathcal{S}_d(n)$ its algebraic matroid and we assume d even, since every skew-symmetric matrix has even rank. The following statement seems to be known, but we have not found an explicit proof of it.

Proposition 3.1. *For even d , $\mathcal{H}_d(n) = \mathcal{S}_d(n)$.*

Proof. Let $d = 2k$. Every skew-symmetric matrix of rank $\leq 2k$ can be expressed as

$$\sum_{l=1}^k (\mathbf{a}_l^T \mathbf{b}_l - \mathbf{b}_l^T \mathbf{a}_l)$$

for some $\mathbf{a}_1, \mathbf{b}_1, \dots, \mathbf{a}_k, \mathbf{b}_k \in \mathbb{R}^n$ (see, e.g., [rgomr10, Lemma 1.3] or [PS05, Theorem 3.23]). Thus, writing $\mathbf{a}_l = (a_{l,1}, \dots, a_{l,n})$ and $\mathbf{b}_l = (b_{l,1}, \dots, b_{l,n})$ we have a surjective map

$$\begin{aligned} T : (\mathbb{R}^n)^{2k} &\rightarrow S_{2k}(n) \subset \mathbb{R}^{\binom{n}{2}} \\ (\mathbf{a}_1, \mathbf{b}_1, \dots, \mathbf{a}_k, \mathbf{b}_k) &\mapsto \sum_{l=1}^k (a_{l,i} b_{l,j} - a_{l,j} b_{l,i})_{1 \leq i < j \leq n}. \end{aligned} \quad (3.1)$$

For each $(\mathbf{a}_l, \mathbf{b}_l)_{l=1,\dots,k} \in (\mathbb{R}^n)^2$, we denote by $J((\mathbf{a}_l, \mathbf{b}_l)_{l=1,\dots,k})$ the Jacobian matrix of T at that point; hence, $\mathcal{S}_{2k}(n)$ equals the linear matroid of rows of $J((\mathbf{a}_l, \mathbf{b}_l)_{l=1,\dots,k})$ for a sufficiently generic choice of $(\mathbf{a}_l, \mathbf{b}_l)_{l=1,\dots,k}$. Now, from Eq. (3.1) it is straightforward to check that $J((\mathbf{a}_l, \mathbf{b}_l)_{l=1,\dots,k})$ is the hyperconnectivity matrix of $(\mathbf{p}_1, \dots, \mathbf{p}_n)$ where

$$\mathbf{p}_i = (b_{1,i}, -a_{1,i}, \dots, b_{k,i}, -a_{k,i}).$$

Thus, $\mathcal{H}_{2k}(n)$ also equals the linear matroid of rows of $J((\mathbf{a}_l, \mathbf{b}_l)_{l=1,\dots,k})$ for a sufficiently generic choice of $(\mathbf{a}_l, \mathbf{b}_l)_{l=1,\dots,k}$. \square

Observe that for $d = 2$ the map T in the above proof is the Plücker embedding of $Gr(2, \mathbb{R})$ into $\mathbb{R}^{\binom{n}{2}}$. Hence, $\mathcal{H}_2(n) = \mathcal{P}_2(n)$ is also the algebraic matroid of $Gr(2, \mathbb{R})$.

4. The generic case

4.1. The three generic rigidity matroids

Let $\mathcal{R}_d(n)$, $\mathcal{H}_d(n)$ and $\mathcal{C}_{d-1}^{d-2}(n)$ denote the *generic* bar-and-joint, hyperconnectivity and cofactor rigidity matroids of dimension d for n points; that is, the matroids $\mathcal{R}_d(\mathbf{p})$, $\mathcal{H}_d(\mathbf{p})$ and $\mathcal{C}_{d-1}^{d-2}(\mathbf{q})$ for sufficiently generic choices of $\mathbf{p} \in (\mathbb{R}^d)^n$ and $\mathbf{q} \in (\mathbb{R}^2)^n$. There are the following conjectures relating them:

Conjecture 4.1 ([Whi96, Conjecture 11.5.1]). $\mathcal{C}_{d-1}^{d-2}(n)$ is the freest matroid in which every K_{d+2} is a circuit. In particular, it is the freest among all abstract rigidity matroids of dimension d .

Conjecture 4.2 ([JT21, Conjecture 6.6(a)]). $\mathcal{H}_d(n)$ is the freest matroid in which every K_{d+2} and every $K_{d+1,d+1}$ are circuits.

Here, we say that \mathcal{N} is *freer than* \mathcal{M} and write $\mathcal{M} \leq \mathcal{N}$, where \mathcal{N} and \mathcal{M} are two matroids on the same ground, if every independent set of \mathcal{M} is also independent in \mathcal{N} . Equivalently, if the rank function of \mathcal{M} is bounded above by that of \mathcal{N} .

Of course, Conjecture 4.1 would imply that $\mathcal{C}_{d-1}^{d-2}(n)$ is freer than both of $\mathcal{R}_d(n)$ and $\mathcal{H}_d(n)$. In turn, a positive answer to [Kal02, Problem 3] would imply $\mathcal{R}_d(n)$ is freer than $\mathcal{H}_d(n)$ (see [Kal02, Section 2.7] for the reason of the implication), so we consider this a conjecture as well:

Conjecture 4.3 (Kalai and Whiteley). For every n and d ,

$$\mathcal{H}_d(n) \leq \mathcal{R}_d(n) \leq \mathcal{C}_{d-1}^{d-2}(n).$$

All these conjectures are trivial for $d = 1$: the three matroids coincide with the graphic matroid, which is the only rigidity matroid in dimension one and satisfies that $K_{2,2}$ is a circuit.

For $d = 2$, Conjecture 4.2 is open (see partial results in [JT21, Section 6.3.1]), but the rest hold true. Indeed, Conjecture 4.3 follows from the more precise statement

$$\mathcal{H}_2(n) < \mathcal{R}_2(n) = \mathcal{C}_1^0(n).$$

The equality on the right is trivial from the definitions, and the strict inequality on the left follows from combining Laman's characterization (which gives $\mathcal{H}_2(n) \leq \mathcal{R}_2(n)$) with the fact that $K_{d+1,d+1}$ is a circuit in $\mathcal{H}_d(n)$ for every d while $K_{3,3}$ is a basis in $\mathcal{R}_2(n)$. The Laman characterization implies also Conjecture 4.1.

In dimension 3 or higher only the following is known:

- Conjecture 4.1 has recently been proven in dimension 3 by Clinch, Jackson and Tanigawa [CJT22], but is open in higher dimensions.
- That result implies $\mathcal{R}_3(n) \leq \mathcal{C}_2^1(n)$, and it is conjectured that $\mathcal{R}_3(n) = \mathcal{C}_2^1(n)$ [Whi96, Conjecture 10.3.2]. In higher dimension we know that $\mathcal{R}_d(n) \neq \mathcal{C}_{d-1}^{d-2}(n)$, since the $(d-4)$ -th cone over $K_{6,7}$ is a basis in $\mathcal{C}_{d-1}^{d-2}(n)$ and dependent in $\mathcal{R}_d(n)$ [Whi96, Section 11], but $\mathcal{R}_d(n) \leq \mathcal{C}_{d-1}^{d-2}(n)$ is open.
- Whether $\mathcal{H}_d(n)$ is less free than $\mathcal{R}_d(n)$ and/or $\mathcal{C}_{d-1}^{d-2}(n)$ is open even for $d = 3$, but we do know graphs that are dependent in $\mathcal{H}_d(n)$ and independent in the other two. For example, no bipartite graph can be spanning in $\mathcal{H}_d(n)$ (see Corollary 4.9 below, which essentially follows [Kal85, Theorem 6.1]); hence, $K_{d+1, \binom{d+1}{2}}$ is dependent in it, while it is a basis in $\mathcal{R}_d(n)$ and $\mathcal{C}_{d-1}^{d-2}(n)$ [Whi96, Sect. 11].

A technique similar to the proof of Theorem 1.1 gives us the following result, supporting the conjecture that $\mathcal{H}_d(n) < \mathcal{R}_d(n)$ for all $d \geq 3$:

Theorem 4.4. *Let $V = X \dot{\cup} Y$ be a vertex set with a given bipartition and let $\{\mathbf{p}_i : i \in V\} \subset \mathbb{R}^{d-1}$ be positions for the vertices. Then, restricted to bipartite graphs with that bipartition the following two d -dimensional rigidity matroids coincide:*

1. *the bar-and-joint matroid of the points $\{(\mathbf{p}_i, 0) : i \in X\} \cup \{(\mathbf{p}_j, 1) : j \in Y\}$.*
2. *the hyperconnectivity matroid of the points $\{(\mathbf{p}_i, 1) : i \in V\}$.*

Proof. Let $E = X \times Y$ be the edge set of the complete bipartite graph $K_{X \times Y}$.

For each point \mathbf{p}_i , choose a linear basis $F^i = (f_1^i, \dots, f_d^i)$ of the space of affine functions $f : \mathbb{R}^{d-1} \rightarrow \mathbb{R}$. For each such choice of bases $\mathcal{F} = (F^i)_{i \in V}$ we define the following *affine rigidity matrix* of $(\mathcal{F}, \mathbf{p})$, of size $|E| \times |V|d$: in the row indexed by the edge $(i, j) \in E$, put the vector $F^i(\mathbf{p}_j)$ in block i , the vector $-F^j(\mathbf{p}_i)$ in block j , and zeroes elsewhere.

With the same argument of Lemma 2.3, the linear matroid of rows of this matrix is independent of the choice of bases. Now, we have that:

- The choice $F^i(\mathbf{x}) = (\mathbf{x} - \mathbf{p}_i, 1)$ for $i \in X$ and $F^j(\mathbf{x}) = (\mathbf{x} - \mathbf{p}_j, -1)$ for $j \in Y$ makes the row of edge (i, j) have the vector $(\mathbf{p}_j - \mathbf{p}_j, 1)$ in block i and the vector $(\mathbf{p}_i - \mathbf{p}_j, -1)$ in block j .
- The choice $F^i(\mathbf{x}) = (\mathbf{x}, 1)$ for every i makes the row of edge (i, j) have the vector $(\mathbf{p}_j, 1)$ in block i and the vector $(-\mathbf{p}_i, -1)$ in block j . \square

Observe that in this proof we only use bipartiteness for the construction in \mathcal{R}_d , but not for \mathcal{H}_d . That is, the *affine rigidity matroid* of the point set $\{\mathbf{p}_1, \dots, \mathbf{p}_n\}$ constructed in the proof coincides with $\mathcal{H}_d(\{(\mathbf{p}_i, 1)\}_{i=1, \dots, n})$ on all graphs, not only bipartite ones.

Corollary 4.5. *Restricted to bipartite graphs we have that $\mathcal{H}_d(n) \leq \mathcal{R}_d(n)$.*

Proof. Consider $\mathcal{H}_d(n)$ realized with generic points $\mathbf{q}_1, \dots, \mathbf{q}_n \in \mathbb{R}^d$. By Lemma 2.1, there is no loss of generality in assuming that they all have last coordinate equal to 1, so we write $\mathbf{q}_i = (\mathbf{p}_i, 1)$ with $\mathbf{p}_i \in \mathbb{R}^{d-1}$. The previous theorem tells us that $\mathcal{H}_d(n) = \mathcal{H}_d(\mathbf{q})$, restricted to bipartite graphs, coincides with \mathcal{R}_d on the point set $\{(\mathbf{p}_i, 0) : i \in X\} \cup \{(\mathbf{p}_j, 1) : j \in Y\}$, which is less free than the generic $\mathcal{R}_d(n)$. \square

4.2. Where does $\mathcal{P}_d(n)$ lie.

We can similarly consider the *generic polynomial rigidity matroid of degree $d - 1$* , that we denote $\mathcal{P}_d(n)$, defined as the matroid $\mathcal{P}_d(t_1, \dots, t_n)$ for a sufficiently generic choice of (t_1, \dots, t_n) . It is obvious that $\mathcal{H}_d(n)$, $\mathcal{R}_d(n)$ and $\mathcal{C}_{d-1}^{d-2}(n)$ are freer than $\mathcal{P}_d(n)$, and the last two strictly so since $K_{\binom{d+1}{2}, d+1}$ is a basis in both \mathcal{R}_d and \mathcal{C}_{d-1}^{d-2} , but it is dependent in \mathcal{P}_d for $d \geq 2$ by Corollary 4.9 below. But we do not know of any graph that is independent in $\mathcal{H}_d(n)$ and dependent in $\mathcal{P}_d(n)$. It is thus plausible that the answer to the next question is positive:

Question 4.6. Is $\mathcal{P}_d(n) = \mathcal{H}_d(n)$ for every n and d ?

In this section we show some common properties of $\mathcal{P}_d(n)$ and $\mathcal{H}_d(n)$. Our first result is that every complete bipartite graph has the same rank in both. It follows from the following lemma, which generalizes [Kal85, Theorem 6.1] (Kalai considers only generic positions).

Lemma 4.7. *Let $G = K_{n_1, n_2}$ be a complete bipartite graph, call $n = n_1 + n_2$ and let $\mathbf{p}_1, \dots, \mathbf{p}_n \in \mathbb{R}^d$ be positions in linear general position (no d of them lie in a linear hyperplane). Then, the rank of G in $\mathcal{H}_d(\mathbf{p})$ equals:*

1. $n_1 n_2$ (that is, G is independent) if $\min\{n_1, n_2\} \leq d$.
2. $dn - d^2$ if $\min\{n_1, n_2\} \geq d$.

Proof. In part (1) we have that all edges have an endpoint of degree at most d , and edges incident to an endpoint of such low degree are automatically independent of the rest as long as they are linearly independent among them, which in our case is guaranteed by general position.

For part (2), an independent subgraph of G with $dn - d^2$ edges can be obtained starting with a $K_{d,d}$ (independent by part (1)) and adding the rest of vertices one by one, connecting each new vertex to d previously added vertices.

As the rigidity matrix has $n_1 n_2$ rows, to finish the proof we only need to show that the orthogonal complement of the column space has dimension at least $(n_1 - d)(n_2 - d)$. We consider columns as living in $\mathbb{R}^{n_1} \otimes \mathbb{R}^{n_2}$ in the natural way, with the entry in row (i, j) representing the coordinate $e_i \otimes e_j$. Regarded in this way, every tensor product $l \otimes m$ of a linear dependence l among the points in one part of $G(\mathbf{p})$ and a linear dependence m among the points

in the other part lies in the orthogonal complement of the columns. By general position, we have $n_1 - d$ and $n_2 - d$ dimensions to choose l and m from, giving a tensor product space of dimension $(n_1 - d)(n_2 - d)$. \square

Remark 4.8. Observe that in part (1) general position is equivalent to “ \mathbf{p} is affinely independent”, and in the proof of part (2) we only use that each part of (G, \mathbf{p}) linearly spans \mathbb{R}^d , a condition weaker than general position. In fact, the same proof gives the more general result: the rank of K_{n_1, n_2} in $\mathcal{H}_d(\mathbf{p})$ equals $n_1 n_2 - (n_1 - d_1)(n_2 - d_2)$ where d_1 and d_2 are the dimensions of the linear spans of the points in the two parts of the framework $(K_{n_1, n_2}, \mathbf{p})$.

This result is the analogue in \mathcal{H}_d of the main result of [BR80]. Although [BR80, Theorem 5.9] gives the rank of K_{n_1, n_2} in $\mathcal{R}_d(\mathbf{p})$ for any choice of \mathbf{p} , we here state only the cases analogue to the ones in Lemma 4.7:

1. If $\min\{n_1, n_2\} \leq d$ and each part is in general position (that is, affinely independent) then K_{n_1, n_2} is independent in $\mathcal{R}_d(\mathbf{p})$; its rank equals $n_1 n_2$.
2. If $\min\{n_1, n_2\} \geq d$ and each part affinely spans \mathbb{R}^d then the rank of K_{n_1, n_2} in $\mathcal{R}_d(\mathbf{p})$ equals $nd - \binom{d+1}{2}$ minus the number of linearly independent quadrics containing \mathbf{p} . That is: if no quadric contains \mathbf{p} (for which a necessary condition is $n_1 + n_2 \geq \binom{d+2}{2}$) then K_{n_1, n_2} is spanning in $\mathcal{R}_d(\mathbf{p})$; from there each quadric containing \mathbf{p} lowers rank by one.

The d -dimensional moment curve is contained in $\binom{d}{2}$ quadrics: for each $2 \leq i \leq d$ we have the quadric $x_i = x_1 x_{i-1}$ ($d - 1$ quadrics) and for each $2 \leq i \leq j \leq d - 1$ we have the quadric $x_1 x_{i+j-1} = x_i x_j$ if $i + j \leq d + 1$ or the quadric $x_d x_{i+j-d} = x_i x_j$ if $i + j > d + 1$ ($\binom{d-1}{2}$ quadrics). Hence, the result of Bolker and Roth recovers that the generic rank of K_{n_1, n_2} in $\mathcal{P}_d(n)$ (that is, the rank in $\mathcal{R}_d(\mathbf{p})$ for \mathbf{p} generic along the moment curve) equals

$$nd - \binom{d+1}{2} - \binom{d}{2} = nd - d^2,$$

as stated in Lemma 4.7.

Corollary 4.9. *Every complete bipartite graph has the same rank in $\mathcal{H}_d(n)$ and $\mathcal{P}_d(n)$, given by the formulas in Lemma 4.7. In particular, in these matroids no bipartite graph is spanning and every $K_{d+1, d+1}$ is a circuit.*

For our next result, recall that a *vertex d -split* on a graph G is the following operation: choose a vertex v of degree at least $d - 1$ with its neighbors divided in three parts A , B , and C with $|B| = d - 1$; remove all edges of the form (v, w) with $w \in C$ and insert a new vertex v' with neighbors $B \cup C \cup \{v\}$. Vertex d -splits are known to preserve generic independence in \mathcal{R}_d and \mathcal{C}_{d-1}^{d-2} . The same holds for \mathcal{P}_d and \mathcal{H}_d as we now show. We call *corank* of a subset X of a matroid the difference $|X| - \text{rank}(X)$.

Proposition 4.10. *Corank does not increase under vertex d -split neither in \mathcal{H}_d nor in \mathcal{P}_d . In particular, vertex splits of independent graphs are independent.*

Proof. The standard proof that vertex split preserves independence in \mathcal{R}_d [Whi90] carries over as follows. To fix notation, let $B = \{v_1, \dots, v_{d-1}\}$ be the neighbors of v that are joined to both v and v' after the split, and let G' be the graph obtained by the split.

Let \mathbf{p} be arbitrary positions for the vertices of G , and denote \mathbf{p}_w the position of each vertex w . Let \mathbf{p}' be positions for G' defined by $\mathbf{p}'_w = \mathbf{p}_w$ for every vertex in G and $\mathbf{p}'_{v'} = \mathbf{p}_v$. Assuming that $\mathbf{p}_v, \mathbf{p}_{v_1}, \dots, \mathbf{p}_{v_{d-1}}$ are linearly independent (that is, a basis of \mathbb{R}^d) we have that every vector orthogonal to all rows of the matrix $H(\mathbf{p})$ becomes orthogonal to all rows of $H(\mathbf{p}')$ by simply repeating in the block of v' what was in the block of v . Conversely, every vector orthogonal to all rows of $H(\mathbf{p}')$ has the same content in the blocks of v and v' . Indeed, if we call $u, u' \in \mathbb{R}^d$ what the vector has in the blocks of v and v' , the vertex split construction implies that u and u' have the same scalar product with the vectors $\mathbf{p}_v, \mathbf{p}_{v_1}, \dots, \mathbf{p}_{v_{d-1}}$.

That is, as long as $\mathbf{p}_v, \mathbf{p}_{v_1}, \dots, \mathbf{p}_{v_{d-1}}$ are a basis, the coranks of G in $\mathcal{H}_d(\mathbf{p})$ and of G' in $\mathcal{H}_d(\mathbf{p}')$ coincide. In particular, this holds for generic positions (matroid $\mathcal{H}_d(n)$) and for generic positions along the moment curve (matroid $\mathcal{P}_d(n)$). Of course, the positions \mathbf{p}' are not generic since $\mathbf{p}'_v = \mathbf{p}'_{v'}$, but perturbing \mathbf{p}' to be generic can only increase the rank of G' and decrease its corank. \square

The following is an example where a vertex split makes the corank decrease in \mathcal{H}_2 : Let G be the graph with edges $\{12, 23, 34, 14, 15, 25, 35, 46, 56\}$. It is a basis (in any abstract rigidity matroid of dimension 2) since vertices 6, 4 and 3, deleted in this order, remove two edges each and result in a K_3 . Let G' be the vertex split of 5 in G with $A = \{1, 3\}$, $B = \{6\}$ and $C = \{2\}$. This is again a basis by the proposition above. It turns out that G' is isomorphic to the vertex split of any vertex in $K_{3,3}$ with $|A| = |B| = |C| = 1$; since $K_{3,3}$ is a circuit, this vertex split decreases the corank from one to zero.

A variation of vertex split is what we here call a *diamond d -split*: choose again a vertex v , now of degree at least d and with its neighbors divided in three parts A, B , and C with $|B| = d$. Then remove all edges of the form (v, w) with $w \in C$ and insert a new vertex v' with neighbors $B \cup C$. Observe that we now do not insert the edge vv' . The following result is proved in [KNN16, Lemma 3.8] in a context that includes \mathcal{H}_d , but restricted to bipartite graphs:

Proposition 4.11. *Corank does not increase under diamond d -split neither in \mathcal{H}_d nor in \mathcal{P}_d . In particular, diamond splits of independent graphs are independent.*

Proof. The proof is very similar to the previous one, and left to the reader. The only significant difference is that now we call $B = \{v_1, \dots, v_d\}$ and assume $\{\mathbf{p}_u : u \in B\}$ to be a linear basis, instead of $\{\mathbf{p}_u : u \in B \cup \{v\}\}$. \square

Let us finish by summing up what these results give as for $d = 2$, in which $\mathcal{P}_2 = \mathcal{H}_2$. In this case all bases of \mathcal{H}_2 are Laman graphs, and they include all *planar* Laman graphs (since every planar Laman graph can be obtained from K_3 by vertex splits [FJW04]), but no *bipartite* Laman graph (since bipartite graphs cannot be rigid in \mathcal{H}_d for any $d \geq 2$). Theorem 6.9 in [JT21], taking into account that independence in \mathcal{H}_2 is preserved under 0-extensions and diamond splits, implies that the case $d = 2$ of Conjecture 4.2 holds if, and only if, it has the property that every dependent flat is a union of copies of K_4 and $K_{3,3}$.

Acknowledgements

We thank Daniel Irving Bernstein, Bill Jackson, Gil Kalai, Meera Sitharam and Walter Whiteley for comments on an early version of this paper. In particular, for directing us to the possible relation between $\mathcal{P}_d(n)$ and $\mathcal{H}_d(n)$. We also thank Eran Nevo for clarifying discussions on the relation between rigidity and shifting, and anonymous referees for helpful comments on the first version of this paper.

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